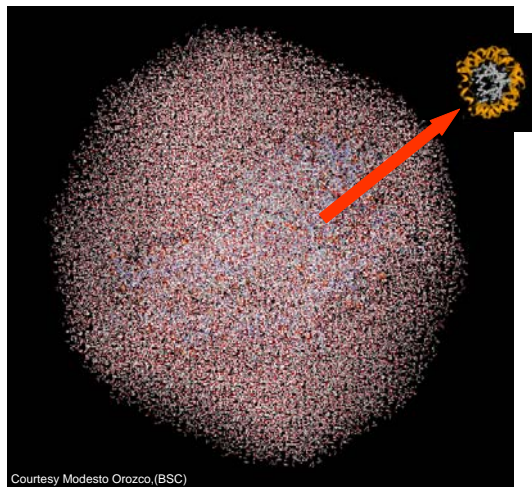


# Advanced Profiling of GROMACS

**Jesus Labarta**  
Director Computer Sciences Research Dept.  
BSC

## All I know about GROMACS

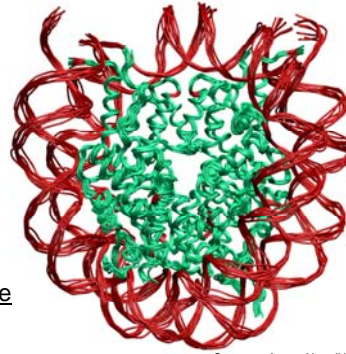
- A Molecular Dynamics application
- Heavily used @ BSC
- Not much ☹



Courtesy Modesto Orozco.(BSC)

# My interest

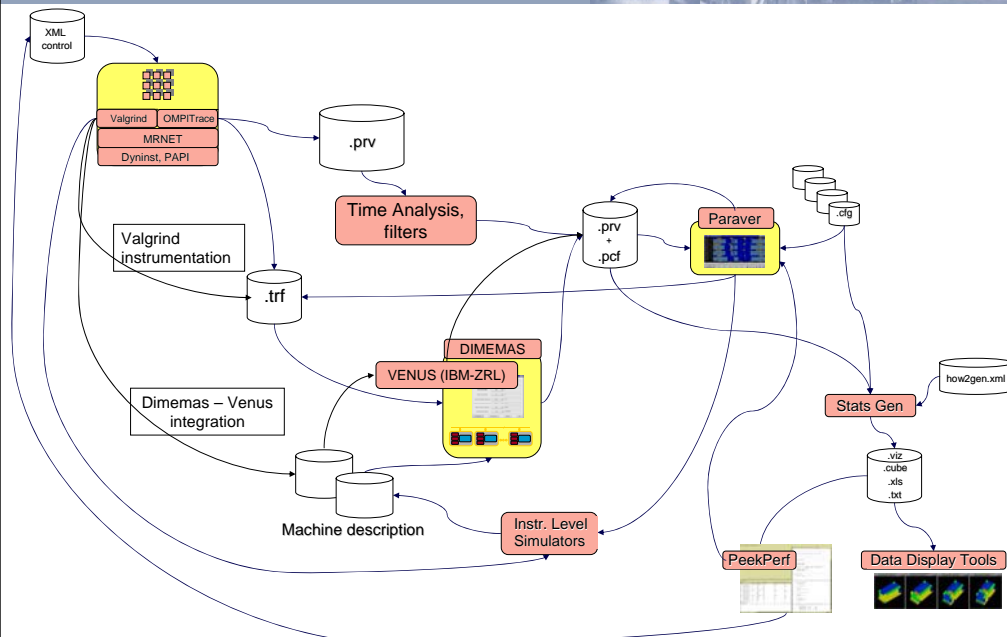
- Efficiency? Scalability?
- Can it be improved?
- Test case for our CEPBA-Tools environment
- Feedback/drive or research
  - Performance analysis tools
  - Programming models
  - Automatic Load balancing
  - Interconnects



Nucleosome test case  
 145732 atoms  
 TIP3P water model  
 parmBSC0 force-field  
 Berendsen thermostat and barostat,  
 2 fs time step  
 constraints on Hydrogens

Courtesy Agnes Noy (UB)

# Evaluation infrastructure



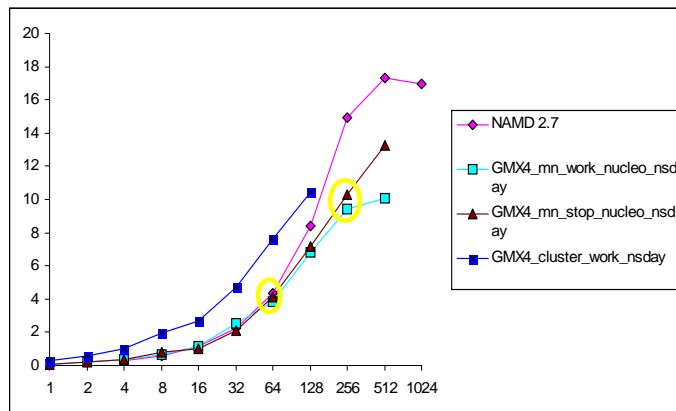
# Index

- A story on
  - Performance
  - for (,;)
    - Parallel efficiency
    - Sequential performance (VMX,...)
  - Hybrid programming



# Performance

- Raw user metric (ns/day)



## Problems?

	Procs	phase	eff
Measured	64	all	0,55
	64	FFTs	0,55
	64	Particles	0,54
	256	all	0,32
	256	Particles	0,35
	256	FFTs	0,29

Not really good

Bad !!!

- Usual suspects
  - Lot of communication, large messages,...
- VMX

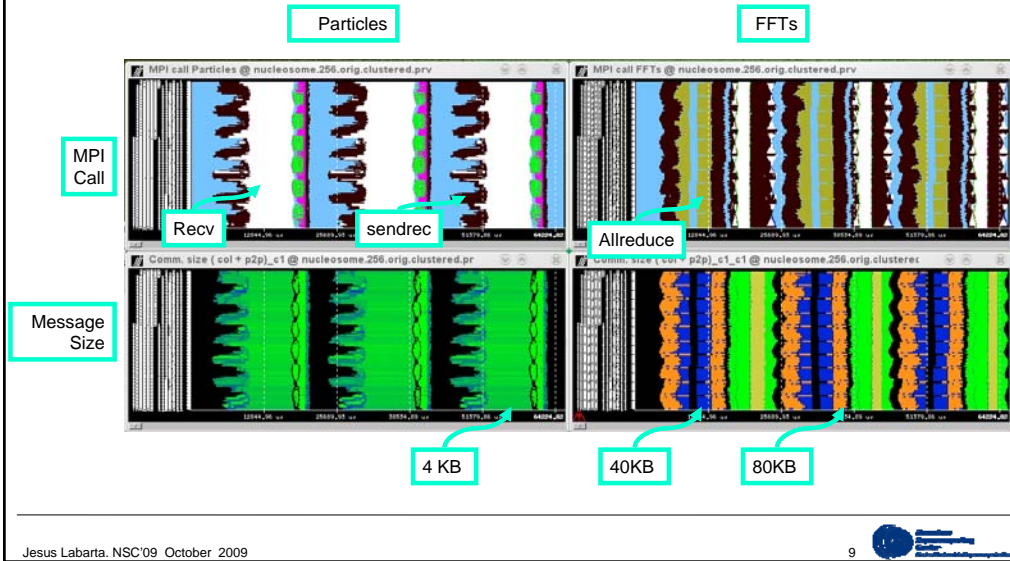


## Parallel efficiency



# Lots of communications, message sizes, ...

- Suspect → convicted?



# Convicted

- Who?
  - Machine? Program?
- Average Bandwidth per core ?
  - 12 -15 MB/s
- Ideal interconnect

	Procs	phase	eff
Measured	64	all	0,55
	64	FFTs	0,55
	64	Particles	0,54
	256	all	0,32
	256	Particles	0,35
Prediction ideal	256	all	0,59
	256	Particles	0,63
	256	FFTs	0,52

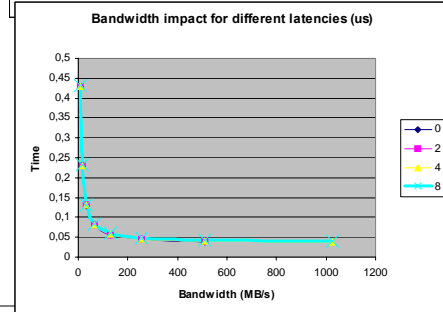
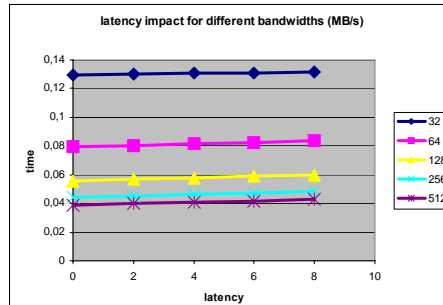
Better !!!!!

Not that good !!!!!

# Sensitivity to interconnect parameters

	Procs	phase	eff
Measured	64	all	0,55
	64	FFTs	0,55
	64	Particles	0,54
	256	all	0,32
	256	Particles	0,35
Prediction ideal	256	all	0,59
	256	Particles	0,63
	256	FFTs	0,52

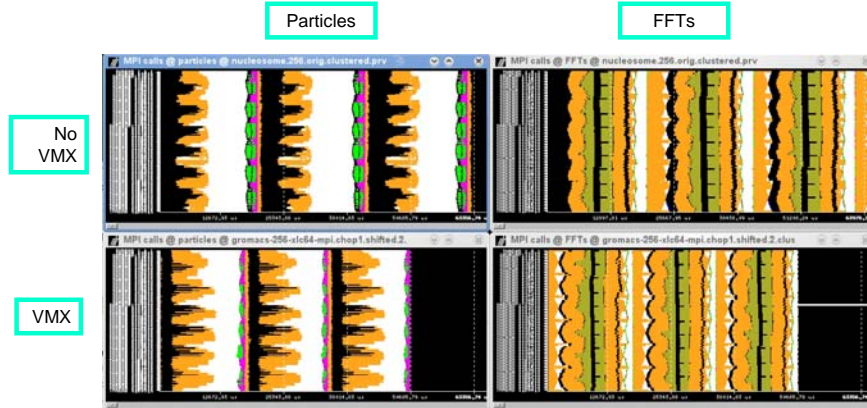
- Latency. Need to group messages?
- Sensitivity to network bandwidth?



# VMX

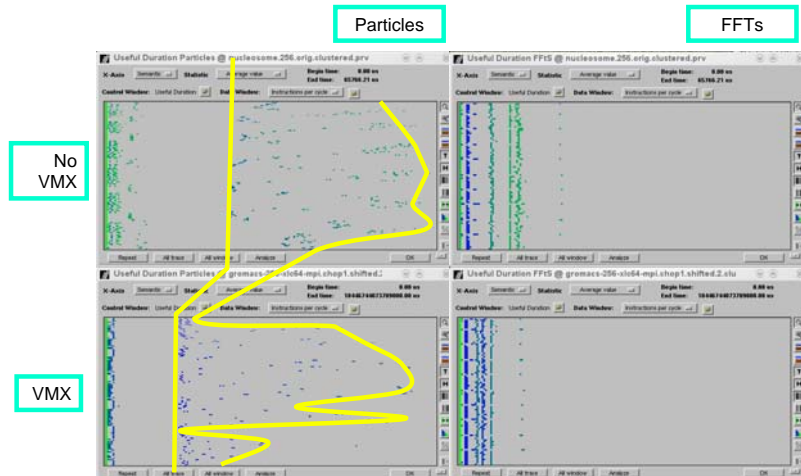
# VMX

- Not as much gain as expected ???



# VMX

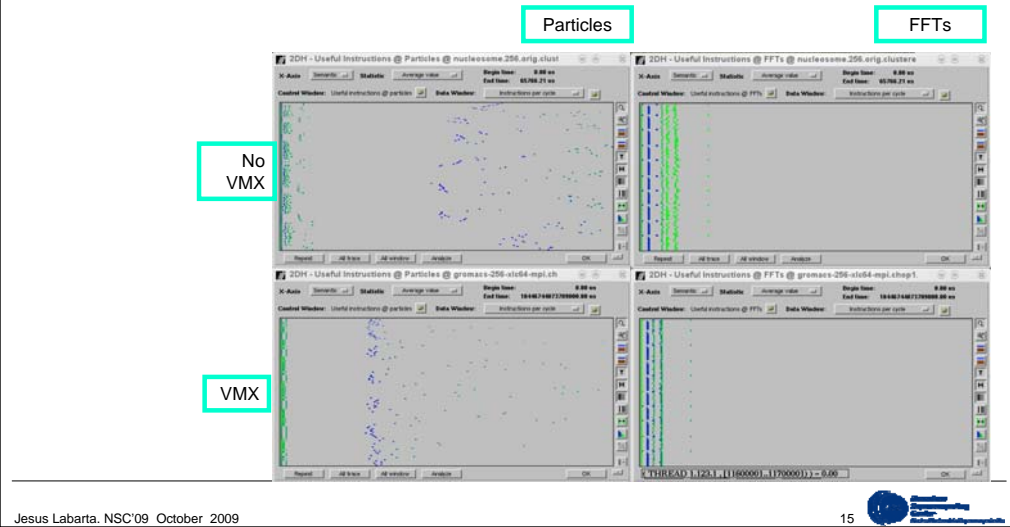
- Histogram of duration: Not uniform impact !!!!
  - Across Particle processes
  - Between FFTs and Particles FFTs





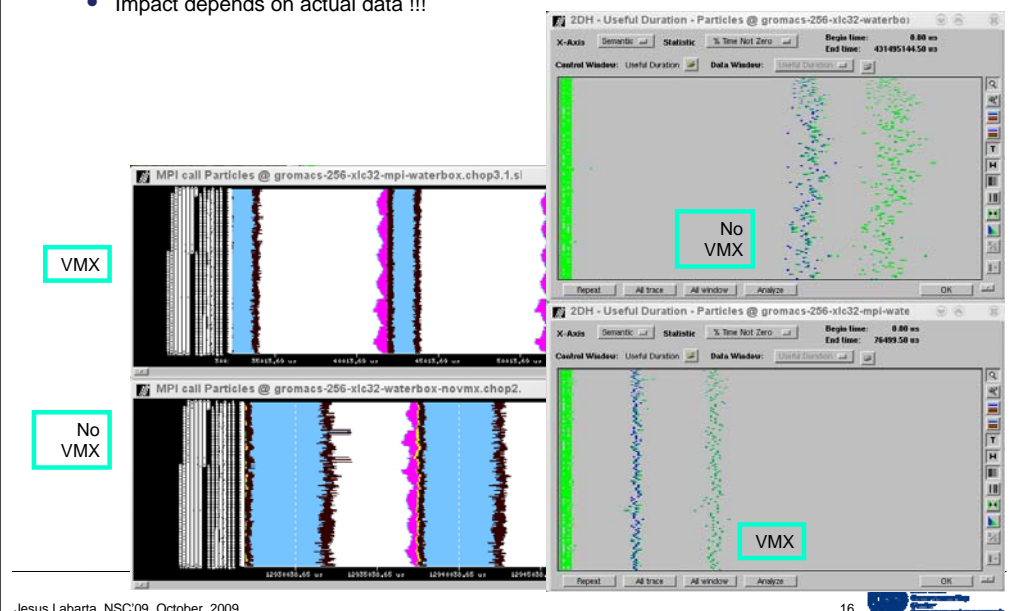
# VMX

- Histograms of Instructions



# VMX @ waterbox

- Impact depends on actual data !!!



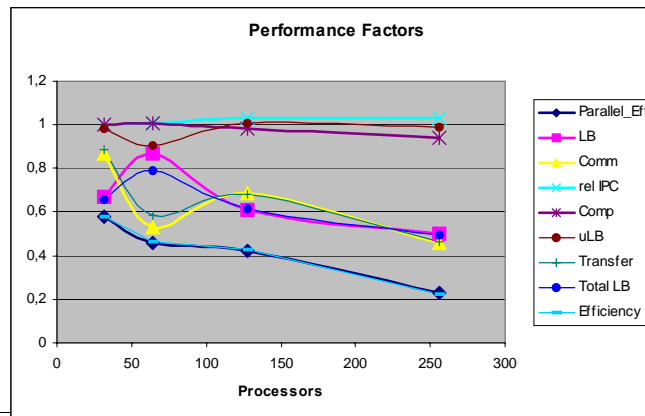


# Back to parallel performance

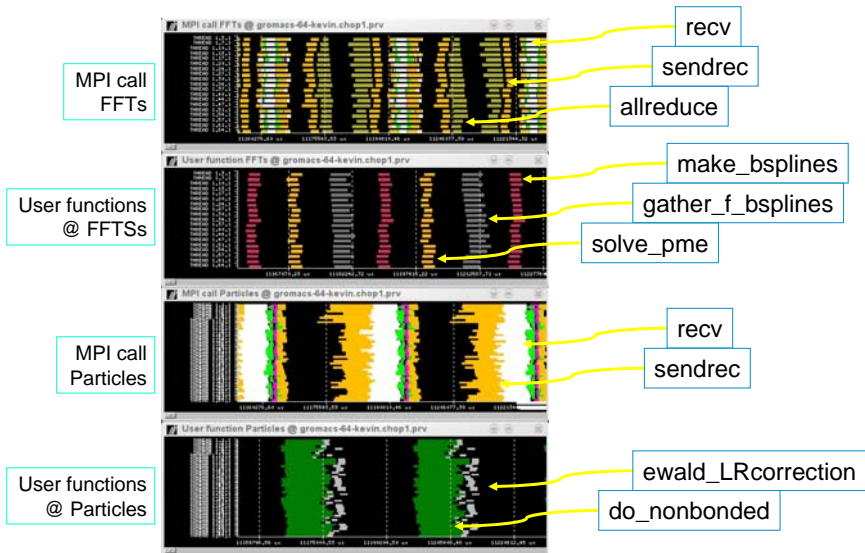
## Performance factors

$$\text{Parallel\_}\eta = \underbrace{LB * \text{microLB} * \text{Transfer}}_{\text{Comm}} * \text{Comp}$$

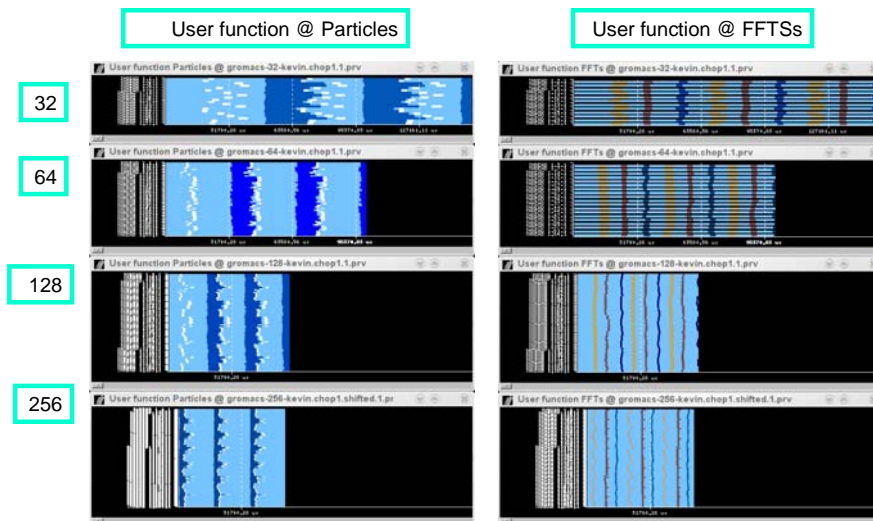
$$\eta = \text{Parallel\_}\eta * \text{relIPC} * \text{Comp}$$



## Structure: A few colors



## Scaling structure

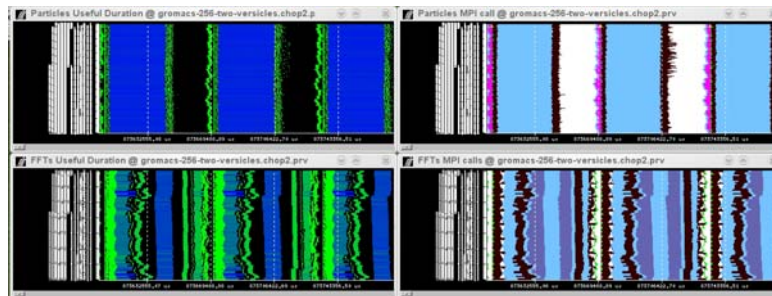


# PRACE Test case

Useful Duration

MPI calls

Particles



FFTs

- Average point to point bandwidth: 13 MB/s

# PRACE Test case

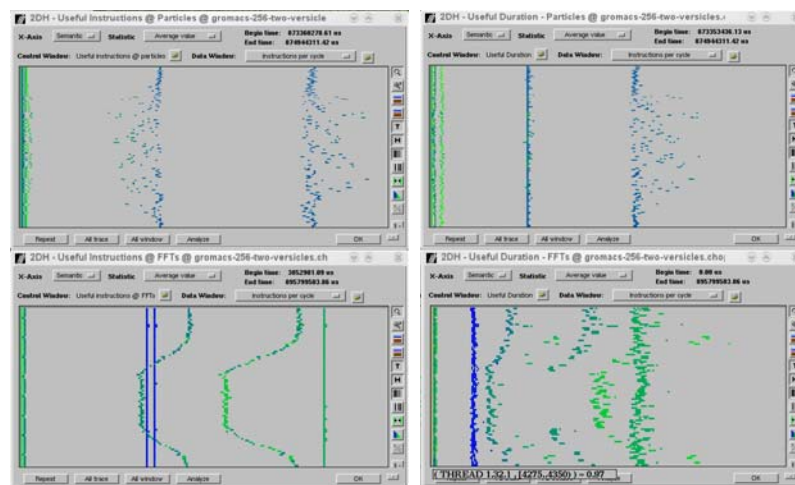
- Effects → opportunities

Instructions

Duration

Particles

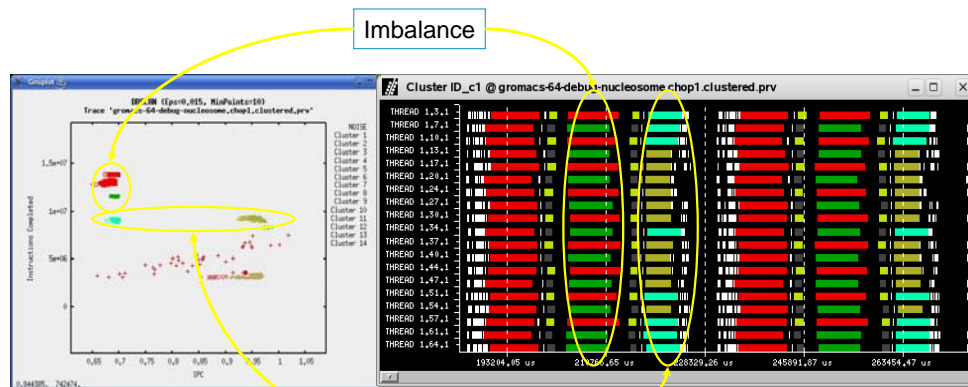
FFTs



# Back to sequential performance

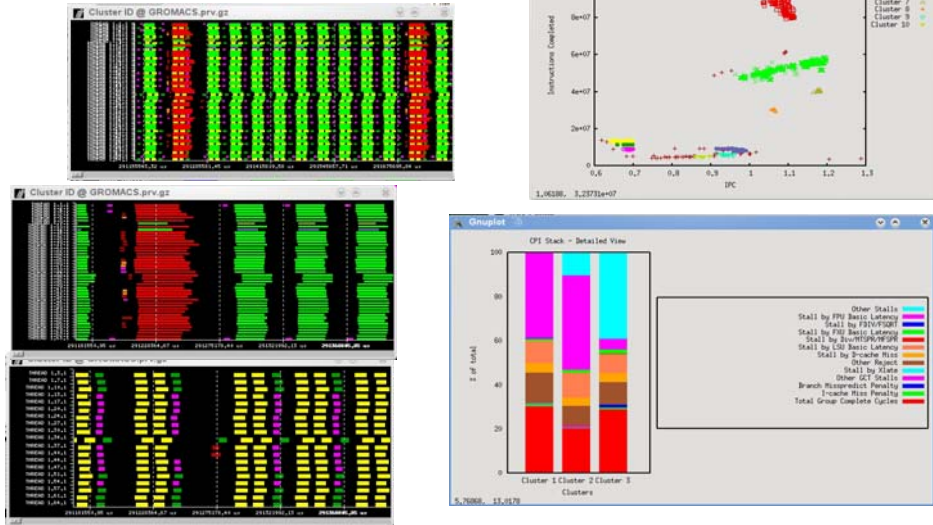
## 64 processes: FFTs

- Globally balanced but some local imbalance



# Online analysis

- Nucleosome

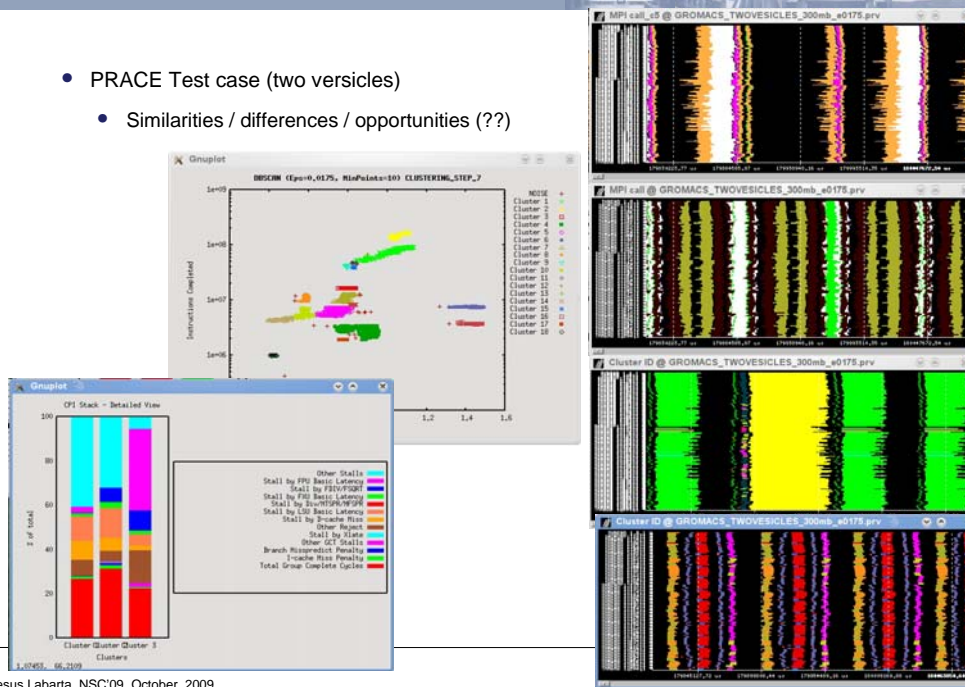


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# Online analysis

- PRACE Test case (two versicles)
  - Similarities / differences / opportunities (??)



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## Detailed metrics

- Application characterization

Metric Description	Cluster 1	Cluster 2	Cluster 3	Cluster 4
<b>FMA ops per floating point instruction</b> PM_FPU_FMA / PM_FPU0_FIN + PM_FPU1_FIN	0,681	0,768	0,700	0,436
<b>Instructions per Load/Store</b> PM_INST_CMPL / (PM_LD_REF_L1 + PM_ST_REF_L1)	1,470	1,465	1,659	1,306
<b>HW floating point instructions (flips)</b> PM_FPU0_FIN + PM_FPU1_FIN	1.174.346	2.896.665	3.549.551	1.465.935
<b>Total floating point operations (flops)</b> PM_FPU0_FIN + PM_FPU1_FIN + PM_FPU_FMA - PM_FPU_STF	1.436.155	3.979.120	6.459.096	3.391.547
<b>Total FP Load&amp;Store operations (fp_tot_ls)</b> PM_LSU_LDF + PM_FPU_STF	1.348.179	3.006.794	3.656.377	2.014.815
<b>FMA %</b> 2 * PM_FPU_FMA / flops	36,51%	44,12%	37,19%	81,21%
<b>Memory Mix</b> (PM_LD_REF_L1+ PM_ST_REF_L1) / PM_INST_DISP	32,61%	24,70%	30,26%	25,20%
<b>Load Mix</b> PM_LD_REF_L1 / PM_INST_DISP	25,83%	19,39%	24,93%	24,76%
<b>Store Mix</b> PM_ST_REF_L1 / PM_INST_DISP	6,77%	5,32%	5,32%	0,44%
<b>FPU Mix</b> flips / PM_INST_DISP	1,06%	8,22%	2,00%	5,32%
<b>FXU Mix</b> PM_FXU_FIN / PM_INST_DISP	20,85%	16,19%	24,38%	16,07%

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## Detailed metrics

- Behavior on architecture

Metric Description	Cluster 1	Cluster 2	Cluster 3	Cluster 4
<b>% Instr. Completed</b> PM_INST_CMPL / PM_INST_DISP	47,93%	36,20%	50,21%	32,90%
<b>L1 misses per Kinstr.</b> (PM_LD_MISS_L1 + PM_ST_MISS_L1 / PM_INST_CMPL) * 1000	10,18	37,04	10,38	63,81
<b>L2 misses per Kinstr.</b> (PM_DATA_FROM_MEM / PM_INST_CMPL) * 1000	0,093	0,103	0,057	1,008
<b>Bytes from main memory per floating point instruction finished</b> (PM_DATA_FROM_MEM * mem_line_size) / PM_FPU_FIN	0,084	0,090	0,068	9,712
<b>Number of Loads per Load miss</b> PM_LD_REF_L1 / PM_LD_MISS_L1	121,81	74,99	102,29	11,84
<b>Number of Stores per Store miss</b> PM_ST_REF_L1 / PM_ST_MISS_L1	24,53	4,91	19,19	50,31
<b>Number of Loads&amp;Stores per L1 miss</b> (PM_LD_REF_L1 + PM_ST_REF_L1) / (PM_LD_MISS_L1 + PM_ST_MISS_L1)	66,80	18,42	58,05	12,00
<b>L1 cache hit rate</b> 1 - (PM_LD_MISS_L1 + PM_ST_MISS_L1) / (PM_LD_REF_L1 + PM_ST_REF_L1)	98,50%	94,57%	98,28%	91,67%
<b>Number of Loads per (D)TLB miss</b> PM_LD_REF_L1 / PM_DTLB_MISS	149.034,25	112.018,98	128.501,34	49.827,87
<b>Number of Loads&amp;Stores per (D)TLB miss</b> (PM_LD_REF_L1 + PM_ST_REF_L1) / PM_DTLB_MISS	188.100,04	142.754,08	155.941,21	50.707,36
<b>% TLB misses per cycle</b> PM_DTLB_MISS / PM_CYC	0,000%	0,000%	0,000%	0,001%
<b>Total Loads from local L2 (M) (total_ld_l2)</b> PM_DATA_FROM_L2 / 1024*1024	0,173	0,060	#VALUE!	0,381
<b>Local L2 load traffic (MB)</b> L1_cache_line_size * total_ld_l2	22,126	7,635	#VALUE!	48,781

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## Detailed metrics

- Real Performance

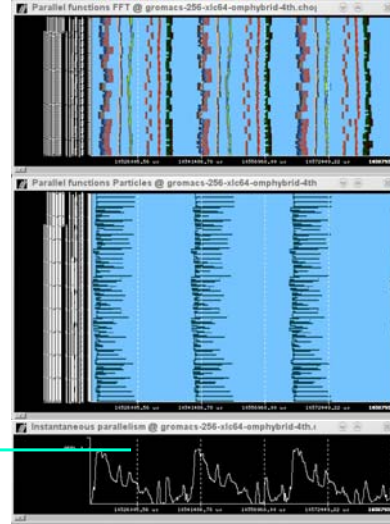
Metric Description	Cluster 1	Cluster 2	Cluster 3	Cluster 4
% Duration	0.60741	0.20625	0.09719	0.04445
Avg. Burst Duration (µs)	21.260,58	8.243,17	36.604,90	5.790,58
Total preempted time (µs)				
Avg. Burst Duration – (PM_CYC * processor_cycle_time)	366,81	123,88	576,02	75,37
% preempted time				
Preempted Time / Avg. Burst Time	1,725%	1,503%	1,574%	1,302%
IPC				
PM_INST_CMPL / PM_CYC	1,11	0,68	1,08	0,69
CPI				
1/IPC	0,90	1,46	0,93	1,45
MIPs				
PM_INST_CMPL / Avg. Burst Duration	2496,91	1547,98	2431,90	1566,34
Mem.BW (MB/s)				
(PM_DATA_FROM_MEM * mem_line_size) / Avg. Burst Duration	29,74	20,43	17,60	202,08
Memory instructions per second				
(PM_LD_REF_L1 + PM_ST_REF_L1) / Avg. Burst Duration	1698,69	1056,39	1465,48	1199,69
HW floating point instructions per cycle				
flops / PM_CYC	0,024	0,155	0,043	0,112
Flop rate (MFLOPs)				
flops / Avg. Burst Duration	67,55	482,72	176,45	585,70
HW floating point instructions rate				
flops / Avg. Burst Duration	55,24	351,40	96,97	253,16
Computation intensity				
flops / fp_tot_Is	1,065	1,323	1,767	1,683
Local L2 load bandwidth per processor (MB/s)				
L1_cache_line_size * total_ld_l2 / Avg. Burst Duration	1.040,72	926,19	#VALUE!	8.424,28
% Loads from local L2 per cycle				
PM_DATA_FROM_L2 / PM_CYC	0,378%	0,335%	#VALUE!	3,043%

## Hybrid MPI + OpenMP



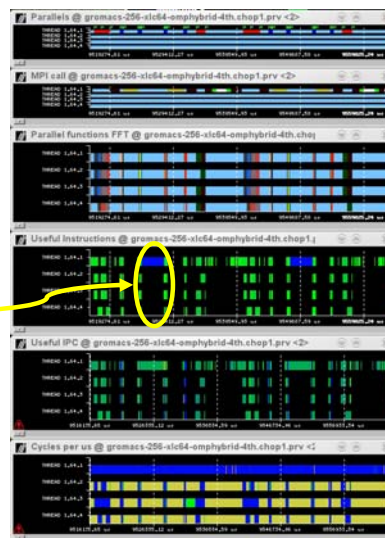
# Hybrid parallelization

- OpenMP parallelization
  - Particles by Christian Simarro (BSC)
  - FFTs by Sebastian von Alfthan (CSC)
- 256 MPI x 4 openMP



# Hybrid parallelization

- FFTs
  - Partial parallelization
  - Too synchronous
    - MPI – OpenMP – MPI – OpenMP ...
  - OS scheduling issues

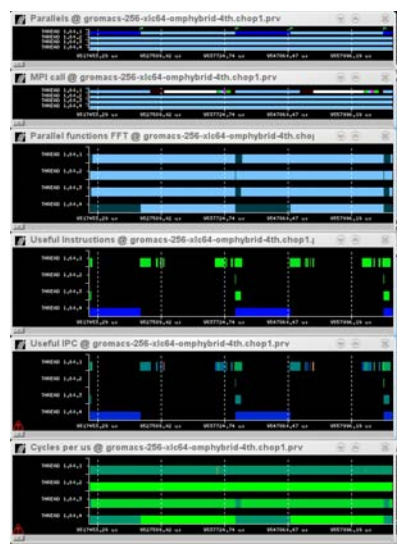


# Hybrid parallelization

- Particles

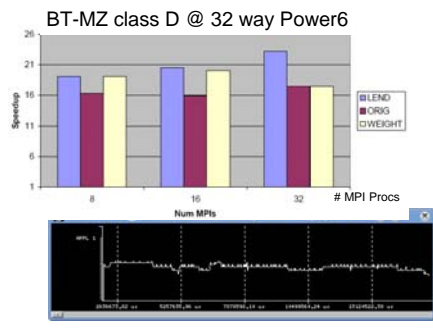
```

for(nen0; nen1);
#pragma omp parallel private(tabledata, nrnb_ind, nlist, outeriter,
for(i=1; i; i++) {
    if (something)
        compute();
}
    
```

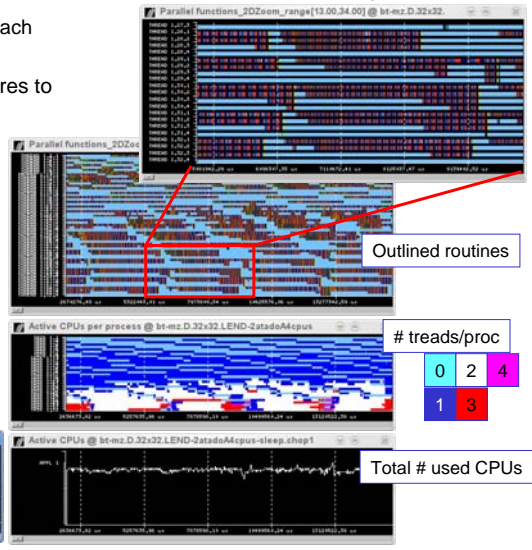


# Dynamic Load Balancing run time

- MPI + OpenMP (StarSs)
  - Job is started with P MPI processes each with 1 OpenMP thread
  - Processes blocking at MPI calls lend cores to other processes
  - GADGET @ 800 cores: Sup=2.5



BT-MZ class D @ 32 way Power6



## Conclusion



- Applications: Nothing IS, everything CAN BE
  - Can always be different, input data dependent
  - Important to get insight into actual behavior.
  - Need dynamic run time optimization
- GROMACS: an interesting case
  - Really heterogeneous application
  - Serialization and load balancing issues
  - A target for MPI + SMPs
    - Automatic Load Balance
    - Overlap communication/computation
- The need for cooperative work

